

# Bethe state preparation

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# Introduction

Quantum state preparation problem:

?

How to prepare a given quantum state (say, an eigenstate of a Hamiltonian) on a quantum computer?

Applications: computing e.g. correlation functions in this state

Attractive candidate: Bethe state (eigenstate of an *integrable* Hamiltonian)

- better understood
- non-trivial

Goal of today's talk: [algorithm for preparing Bethe states of the Heisenberg quantum spin chain](#)

# Outline

1. Bethe ansatz review
2. Dicke state preparation
3. Bethe state preparation
4. Bethe roots from VQE
5. Outlook

# I. Bethe ansatz review



# Coordinate Bethe ansatz (1931)

Remarkable solution!

Reduces the problem to solving a system of polynomial equations “Bethe equations”





Eigenvectors are **multi-particle** (“magnon”) states

ground (0-particle) state:

$$|\psi_0\rangle = \left( \begin{array}{c} 1 \\ 0 \end{array} \right)^{\otimes L} = |\uparrow \cdots \uparrow\rangle$$

$$H|\psi_0\rangle = 0$$

1-particle state:

$$|\psi(k)\rangle = \sum_{x=1}^L e^{ikx} |\uparrow \cdots \underset{x}{\downarrow} \cdots \uparrow\rangle$$

$$H|\psi(k)\rangle = e(k)|\psi(k)\rangle$$

$$e(k) = 4 \sin^2\left(\frac{k}{2}\right)$$

1-particle energy

provided

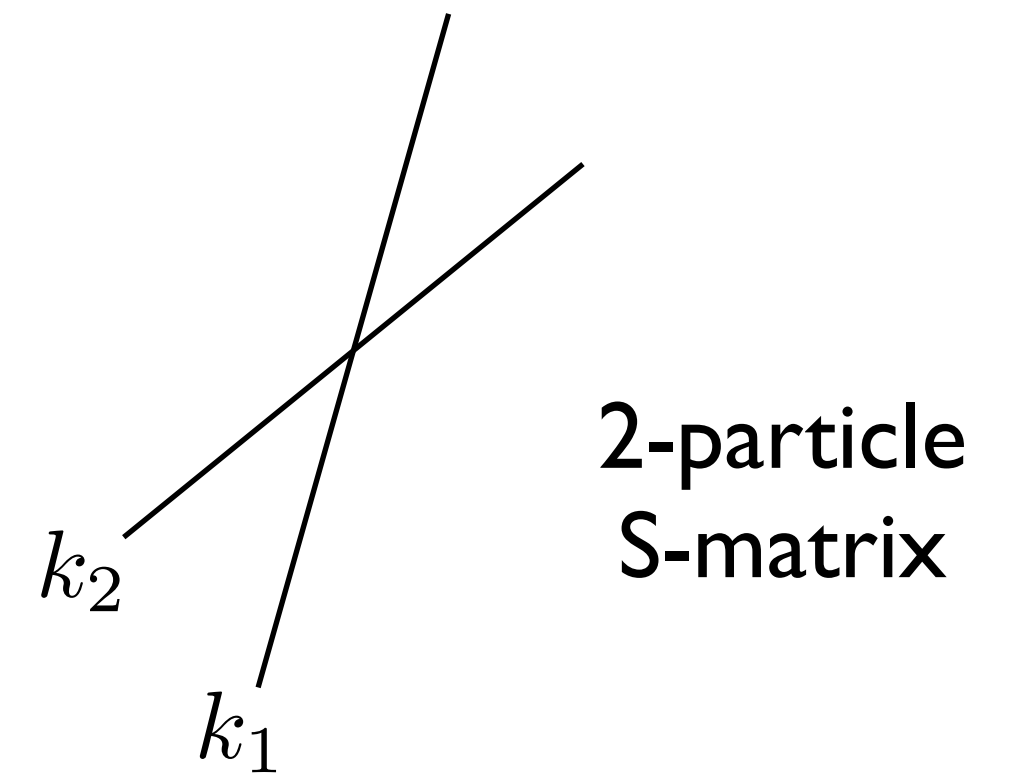
$$e^{ikL} = 1$$

2-particle state:

$$|\psi(k_1, k_2)\rangle = \sum_{1 \leq x_1 < x_2 \leq L} \left[ s(k_2, k_1) e^{i(k_1 x_1 + k_2 x_2)} - s(k_1, k_2) e^{i(k_2 x_1 + k_1 x_2)} \right] |\uparrow \cdots \downarrow_{x_1} \cdots \downarrow_{x_2} \cdots \uparrow\rangle$$

$$s(k, k') = 1 - 2e^{ik'} + e^{i(k+k')}$$

$$S(k_2, k_1) = -\frac{s(k_1, k_2)}{s(k_2, k_1)}$$



$$H|\psi(k_1, k_2)\rangle = E|\psi(k_1, k_2)\rangle \quad E = e(k_1) + e(k_2)$$

provided  $\begin{cases} e^{ik_1 L} = S(k_1, k_2) \\ e^{ik_2 L} = S(k_2, k_1) \end{cases}$

M-particle state:

$$|\psi(k_1, \dots, k_M)\rangle = \sum_{1 \leq x_1 < \dots < x_M \leq L} \sum_{\sigma \in \text{Perm}(1, \dots, M)} \varepsilon(\sigma) A(k_{\sigma(1)}, \dots, k_{\sigma(M)}) e^{i \sum_{j=1}^M k_{\sigma(j)} x_j} | \uparrow \dots \downarrow \dots \downarrow \dots \uparrow \rangle_{x_1 \quad x_M}$$

Bethe state

$$A(k_1, \dots, k_M) = \prod_{1 \leq j < l \leq M} s(k_l, k_j) \quad \varepsilon(\sigma) = \pm 1 \quad \text{signature of } \sigma$$

$$H|\psi(k_1, \dots, k_M)\rangle = E|\psi(k_1, \dots, k_M)\rangle \quad E = \sum_{j=1}^M e(k_j)$$

provided

$$e^{ik_j L} = \prod_{l=1; l \neq j}^M S(k_j, k_l), \quad j = 1, \dots, M$$

Bethe equations

$$\left( \frac{u_j + \frac{i}{2}}{u_j - \frac{i}{2}} \right)^L = \prod_{\substack{l \neq j \\ l=1}}^M \frac{u_j - u_l + i}{u_j - u_l - i}$$

$$u_j = u(k_j)$$

$$u(k) = \frac{1}{2} \cot\left(\frac{k}{2}\right)$$

SU(2) symmetry  $\Rightarrow$  degeneracy  $L - 2M + 1$

$$M \leq \frac{L}{2}$$

Example:  $L = 4$   $M = 0, 1, 2$

$M$	$u_j$	$E$	degeneracy
0	-	0	5
1	$\frac{1}{2}$	2	3
1	$-\frac{1}{2}$	2	3
1	0	4	3
2	$\pm \frac{i}{2}$	2	1
2	$\pm \frac{1}{2\sqrt{3}}$	6	1

total:  $16 = 2^4 = 2^L$  “complete” ✓

## 2. Dicke state preparation

# Dicke states

$$|D_k^n\rangle : \text{completely symmetric state of } |1\rangle\text{'s and } |0\rangle\text{'s}$$

$$\begin{array}{ccc} \uparrow & \uparrow & \\ k & n-k & \end{array} \quad \text{total \# qubits} = n$$

$$|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \quad |1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

Ex:

$$|D_2^4\rangle = \frac{1}{\sqrt{6}} (|1100\rangle + |1010\rangle + |0110\rangle + |1001\rangle + |0101\rangle + |0011\rangle)$$

$$|0011\rangle = |0\rangle \otimes |0\rangle \otimes |1\rangle \otimes |1\rangle$$

$$|D_k^n\rangle \propto (\mathbb{S}^-)^k |0\rangle^{\otimes n} \quad \mathbb{S}^- = \mathbb{S}^x - i\mathbb{S}^y \quad \vec{\mathbb{S}} = \sum_{i=1}^n \vec{S}_i \quad \vec{S}_i = \frac{1}{2}\vec{\sigma}_i$$

Exact ground states of ferromagnetic Heisenberg & Lipkin-Meshkov-Glick Hamiltonians

$$-\sum_i \vec{S}_i \cdot \vec{S}_{i+1} \quad -\vec{\mathbb{S}}^2 = -\sum_{i,j} \vec{S}_i \cdot \vec{S}_j$$

How to prepare on quantum computer?

Cannot implement  $|D_k^n\rangle \propto (\mathbb{S}^-)^k |0\rangle^{\otimes n}$        $\mathbb{S}^-$  is not unitary

Instead:

$$|e_k^n\rangle = |0\rangle^{\otimes(n-k)} |1\rangle^{\otimes k} \quad \text{“reference” state (product)}$$

We seek:

$$U_n |e_k^n\rangle = |D_k^n\rangle$$

“Dicke operator”

- unitary
- independent of  $k$

key idea: recursion!

[Bärtschi, Eidenbenz 2019]

$$\text{Ex: } |D_2^4\rangle = \frac{1}{\sqrt{6}} (|1100\rangle + |1010\rangle + |0110\rangle + |1001\rangle + |0101\rangle + |0011\rangle)$$
$$(\underbrace{|110\rangle + |101\rangle + |011\rangle}_{(|110\rangle + |101\rangle + |011\rangle) \otimes |0\rangle}) \otimes |0\rangle + (\underbrace{|100\rangle + |010\rangle + |001\rangle}_{(|100\rangle + |010\rangle + |001\rangle) \otimes |1\rangle}) \otimes |1\rangle$$

$$|D_2^4\rangle = \sqrt{\frac{1}{2}} |D_2^3\rangle \otimes |0\rangle + \sqrt{\frac{1}{2}} |D_1^3\rangle \otimes |1\rangle$$

$$|D_k^n\rangle = \sqrt{\frac{n-k}{n}} |D_k^{n-1}\rangle \otimes |0\rangle + \sqrt{\frac{k}{n}} |D_{k-1}^{n-1}\rangle \otimes |1\rangle$$

Use  $|D_k^n\rangle = U_n |e_k\rangle$  on both sides:

$$U_n |e_k\rangle = (U_{n-1} \otimes \mathbb{I}) \left( \underbrace{\sqrt{\frac{n-k}{n}} |e_k^{n-1}\rangle \otimes |0\rangle + \sqrt{\frac{k}{n}} |e_{k-1}^{n-1}\rangle \otimes |1\rangle}_{\equiv W_n |e_k^n\rangle} \right)$$

$$\equiv W_n |e_k^n\rangle$$

- unitary
- independent of  $k$



$$W_n |e_k^n\rangle = \sqrt{\frac{n-k}{n}} |e_k^{n-1}\rangle \otimes |0\rangle + \sqrt{\frac{k}{n}} |e_{k-1}^{n-1}\rangle \otimes |1\rangle$$

$$U_n = (U_{n-1} \otimes \mathbb{I}) W_n$$

$\Rightarrow$

$$U_n = \prod_{m=2}^{\overset{\curvearrowright}{n}} \left( W_m \otimes \mathbb{I}^{\otimes(n-m)} \right)$$

Suffices to construct  $W_m$ 's !

## Constructing $W_m$ 's

Strategy: look for operators  $I_{m,l}$  such that

$$I_{m,l'} |e_l^m\rangle = \begin{cases} |e_l^m\rangle & l' < l \\ W_m |e_l^m\rangle & l' = l \end{cases}$$

and

$$I_{m,l'} (I_{m,l} |e_l^m\rangle) = (I_{m,l} |e_l^m\rangle) \quad \text{for } l' > l$$

$$|e_l^m\rangle = |0\rangle^{\otimes(m-l)} |1\rangle^{\otimes l}$$

perform  $W_m |e_l^m\rangle$  for fixed  $l$

ordered so as to not interfere with each other

Then

$$W_m = \prod_{l=1}^{\overset{\curvearrowright}{m-1}} I_{m,l}$$

Suffices to construct  $I_{m,l}$ 's !

## Constructing $I_{m,l}$ 's

$$I_{m,l}|e_l^m\rangle = W_m |e_l^m\rangle$$

$$= W_m |0\rangle^{\otimes(m-l)} |1\rangle^{\otimes l} = \sqrt{\frac{m-l}{m}} |0\rangle^{\otimes(m-l-1)} |1\rangle^{\otimes l} \otimes |0\rangle + \sqrt{\frac{l}{m}} |0\rangle^{\otimes(m-l)} |1\rangle^{\otimes(l-1)} \otimes |1\rangle$$

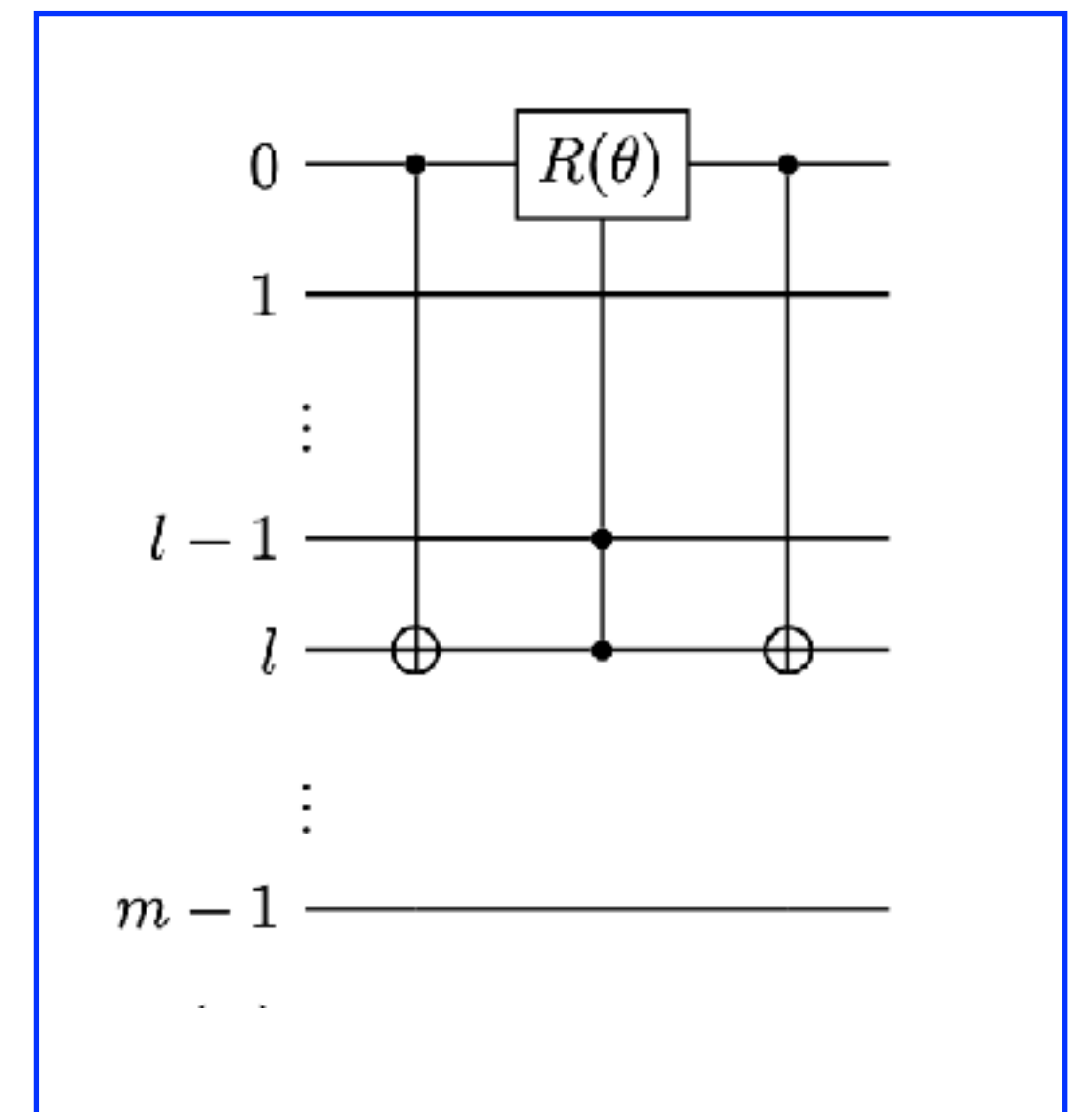
$$\begin{array}{ccc} l & l-1 & 0 \\ \downarrow & \downarrow & \downarrow \\ |0 \dots 0 & 1 \dots 1 & 1\rangle \end{array}$$

$$\begin{array}{ccc} l & l-1 & 0 \\ \downarrow & \downarrow & \downarrow \\ |0 \dots 1 & 1 \dots 1 & 0\rangle \end{array}$$

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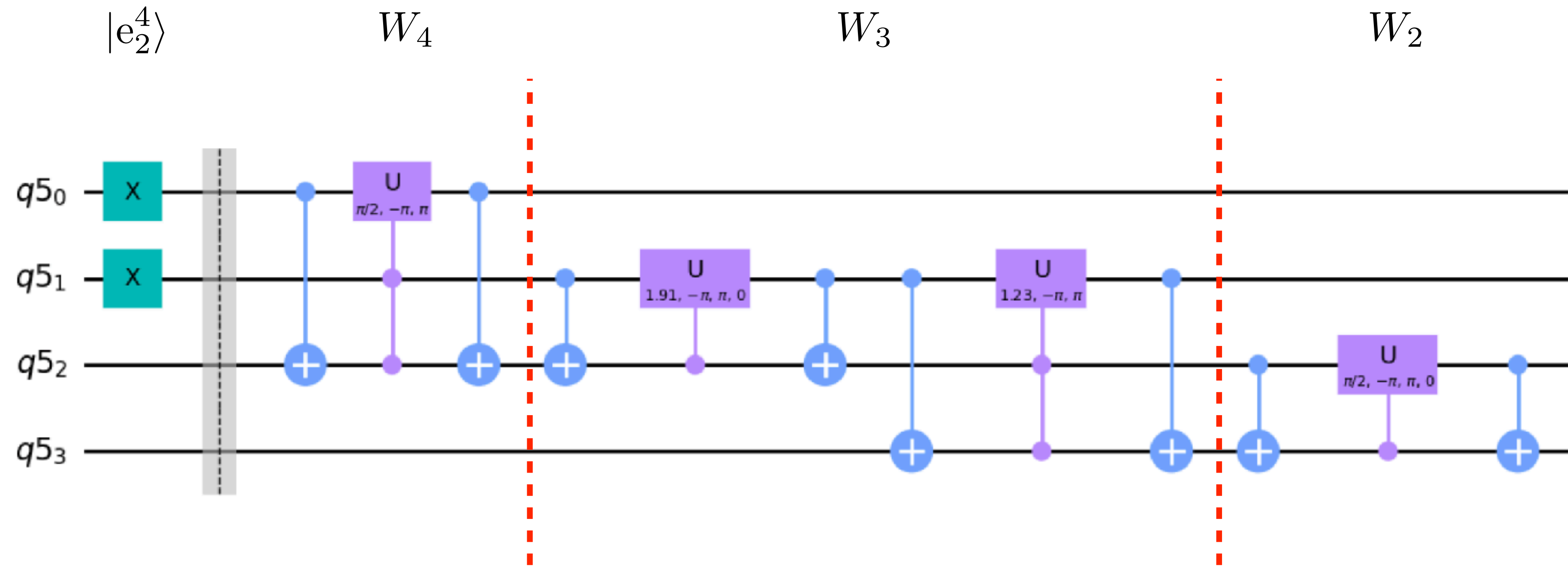
$$I_{m,l} : |0\rangle_l |1\rangle_{l-1} |1\rangle_0 \mapsto \underbrace{\sqrt{\frac{m-l}{m}} |1\rangle_l |1\rangle_{l-1} |0\rangle_0}_{-\sin(\frac{\theta}{2})} + \underbrace{\sqrt{\frac{l}{m}} |0\rangle_l |1\rangle_{l-1} |1\rangle_0}_{\cos(\frac{\theta}{2})}$$

$$R(\theta) = \begin{pmatrix} \cos(\frac{\theta}{2}) & -\sin(\frac{\theta}{2}) \\ \sin(\frac{\theta}{2}) & \cos(\frac{\theta}{2}) \end{pmatrix}$$



circuit size =  $\mathcal{O}(M(L - M))$

Ex:  $|D_2^4\rangle$



## Generalizations:

- higher spin                      RN, Ravanini, Raveh      arXiv: 2402.03233
- higher rank                      RN, Raveh                      arXiv: 2301.04989
- q-deformation                      Raveh, RN                      arXiv: 2308.08392

### 3. Bethe state preparation

Closed periodic spin-1/2 XXZ chain

$$\mathcal{H} = -\frac{1}{2} \sum_{n=1}^L \left( \sigma_n^x \sigma_{n+1}^x + \sigma_n^y \sigma_{n+1}^y + \Delta \left( \sigma_n^z \sigma_{n+1}^z - \mathbb{I} \right) \right), \quad \vec{\sigma}_{L+1} = \vec{\sigma}_1$$

Assume Bethe roots  $\{k_1, \dots, k_M\}$  are known.

How to prepare corresponding Bethe state  $|B_M^L\rangle$  on quantum computer?

Previous:

- Van Dyke, Barron, Mayhall, Barnes, Economou 2021 coordinate BA probabilistic; ancillas, real Bethe roots
- Sopena, Gordon, García-Martín, Sierra, López 2022 algebraic BA deterministic; QR decompositions
- Ruiz, Sopena, Gordon, Sierra, López 2023 algebraic/coordinate BA proposed analytical formulae for unitaries

Today: D. Raveh, RN 2024 coordinate BA deterministic; no ancillas or QR decomp

set of all permutations of M 1's and L-M 0's

Bethe state

$$|B_M^L\rangle \propto \sum_{w \in P(L,M)} f(w) |w\rangle$$

Ex:  $|B_2^4\rangle \propto f(0011) |0011\rangle + f(0101) |0101\rangle + f(1001) |1001\rangle + f(0110) |0110\rangle + f(1010) |1010\rangle + f(1100) |1100\rangle$

Closed periodic spin-1/2 XXZ chain

$$\mathcal{H} = -\frac{1}{2} \sum_{n=1}^L \left( \sigma_n^x \sigma_{n+1}^x + \sigma_n^y \sigma_{n+1}^y + \Delta \left( \sigma_n^z \sigma_{n+1}^z - \mathbb{I} \right) \right), \quad \vec{\sigma}_{L+1} = \vec{\sigma}_1$$

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Bethe state

$$|B_M^L\rangle \propto \sum_{w \in P(L, M)} f(w) |w\rangle$$

$$f(w) = \sum_{\sigma \in \text{Perm}(1, \dots, M)} \varepsilon(\sigma) A(k_{\sigma(1)}, \dots, k_{\sigma(M)}) e^{i \sum_{j=1}^M k_{\sigma(j)} x_j}$$

$x_j \in \{1, \dots, L\}$       positions of 1's in w

Ex:

$$|B_2^4\rangle \propto f(0011) |0011\rangle + f(0101) |0101\rangle + f(1001) |1001\rangle + f(0110) |0110\rangle + f(1010) |1010\rangle + f(1100) |1100\rangle$$



Closed periodic spin-1/2 XXZ chain

$$\mathcal{H} = -\frac{1}{2} \sum_{n=1}^L \left( \sigma_n^x \sigma_{n+1}^x + \sigma_n^y \sigma_{n+1}^y + \Delta \left( \sigma_n^z \sigma_{n+1}^z - \mathbb{I} \right) \right), \quad \vec{\sigma}_{L+1} = \vec{\sigma}_1$$

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Bethe state  $|B_M^L\rangle \propto \sum_{w \in P(L,M)} f(w) |w\rangle$        $f(w) = \sum_{\sigma \in \text{Perm}(1, \dots, M)} \varepsilon(\sigma) A(k_{\sigma(1)}, \dots, k_{\sigma(M)}) e^{i \sum_{j=1}^M k_{\sigma(j)} x_j}$

Ex:  $|B_2^4\rangle \propto f(0011) |0011\rangle + f(0101) |0101\rangle + f(1001) |1001\rangle + f(0110) |0110\rangle + f(1010) |1010\rangle + f(1100) |1100\rangle$        $\binom{L}{M}$  known!

$$|\psi(b)\rangle = \frac{1}{F(b)} \sum_{\substack{w \in P(L,M) \\ w=ab}} f(w) |w\rangle$$

“b-tail” of  $|B_M^L\rangle$  (ends in b)

permutations  $w$  of the form  $a$   $b$   
(concatenation)

$$|w\rangle = |a\rangle|b\rangle$$

**Ex: For**  $|B_2^4\rangle \propto f(0011) |0011\rangle + f(0101) |0101\rangle + f(1001) |1001\rangle$   
 $+ f(0110) |0110\rangle + f(1010) |1010\rangle + f(1100) |1100\rangle$

$$|\psi(0)\rangle = \frac{1}{F(0)} (f(0110) |0110\rangle + f(1010) |1010\rangle + f(1100) |1100\rangle)$$

$$|B_M^L\rangle = |\psi(\{\})\rangle$$

$$b = \{\}$$

We'll do recursion over length of  $b$ !

Define:

$$U_m |0\rangle^{\otimes(m-l)} |1\rangle^{\otimes l} |b\rangle_{L-m} = |\psi(b)\rangle_L$$

- “Bethe operator”
- unitary
  - independent of  $l, b$

$$b = \{ \} :$$

$$|B_M^L\rangle = U_L |0\rangle^{\otimes(L-M)} |1\rangle^{\otimes M}$$

Recursion:

$$|\psi(b)\rangle = G(0b) |\psi(0b)\rangle + G(1b) |\psi(1b)\rangle$$

$$G(ib) = \frac{F(ib)}{F(b)}, \quad i = 0, 1$$

Ex: For  $|\psi(0)\rangle = \frac{1}{F(0)} (f(1100) |1100\rangle + f(0110) |0110\rangle + f(1010) |1010\rangle)$

Define:

$$U_m |0\rangle^{\otimes(m-l)} |1\rangle^{\otimes l} |b\rangle_{L-m} = |\psi(b)\rangle_L$$

↑

$b = \{ \} :$

- “Bethe operator”
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$$|B_M^L\rangle = U_L |0\rangle^{\otimes(L-M)} |1\rangle^{\otimes M}$$

Recursion:

$$|\psi(b)\rangle = G(0b) |\psi(0b)\rangle + G(1b) |\psi(1b)\rangle$$

$$G(ib) = \frac{F(ib)}{F(b)}, \quad i = 0, 1$$

Ex: For

$$|\psi(0)\rangle = \frac{1}{F(0)} (f(1100) |1100\rangle + f(0110) |0110\rangle + f(1010) |1010\rangle)$$

$$|\psi(0)\rangle = \frac{G(00)}{F(00)} f(1100) |1100\rangle + \frac{G(10)}{F(10)} (f(0110) |0110\rangle + f(1010) |1010\rangle)$$

$\underbrace{\hspace{10em}}_{= |\psi(00)\rangle} \qquad \qquad \qquad \underbrace{\hspace{10em}}_{= |\psi(10)\rangle}$

$$|\psi(0)\rangle = G(00) |\psi(00)\rangle + G(10) |\psi(10)\rangle \quad \checkmark$$

Define:

$$U_m |0\rangle^{\otimes(m-l)} |1\rangle^{\otimes l} |b\rangle_{L-m} = |\psi(b)\rangle_L \quad (1)$$

$b = \{ \} :$

“Bethe operator”

- unitary
- independent of  $l, b$

$$|B_M^L\rangle = U_L |0\rangle^{\otimes(L-M)} |1\rangle^{\otimes M}$$

Recursion:

$$|\psi(b)\rangle = G(0b) |\psi(0b)\rangle + G(1b) |\psi(1b)\rangle \quad (2)$$

$$G(ib) = \frac{F(ib)}{F(b)}, \quad i = 0, 1$$

Use (1) on both sides of (2) to obtain recursion for  $U_m$

$$U_m = U_{m-1} W_m$$

$\Rightarrow$

$$U_L = \prod_{m=2}^{\overset{\curvearrowright}{L}} W_m$$

$$W_m |0\rangle^{\otimes(m-l)} |1\rangle^{\otimes l} |b\rangle_{L-m} = G(0b) |0\rangle^{\otimes(m-l-1)} |1\rangle |1\rangle^{\otimes(l-1)} |0\rangle |b\rangle_{L-m} \\ + G(1b) |0\rangle^{\otimes(m-l-1)} |0\rangle |1\rangle^{\otimes(l-1)} |1\rangle |b\rangle_{L-m}$$

Suffices to construct  $W_m$ 's !

Strategy: look for operators  $I_{m,l}$  such that

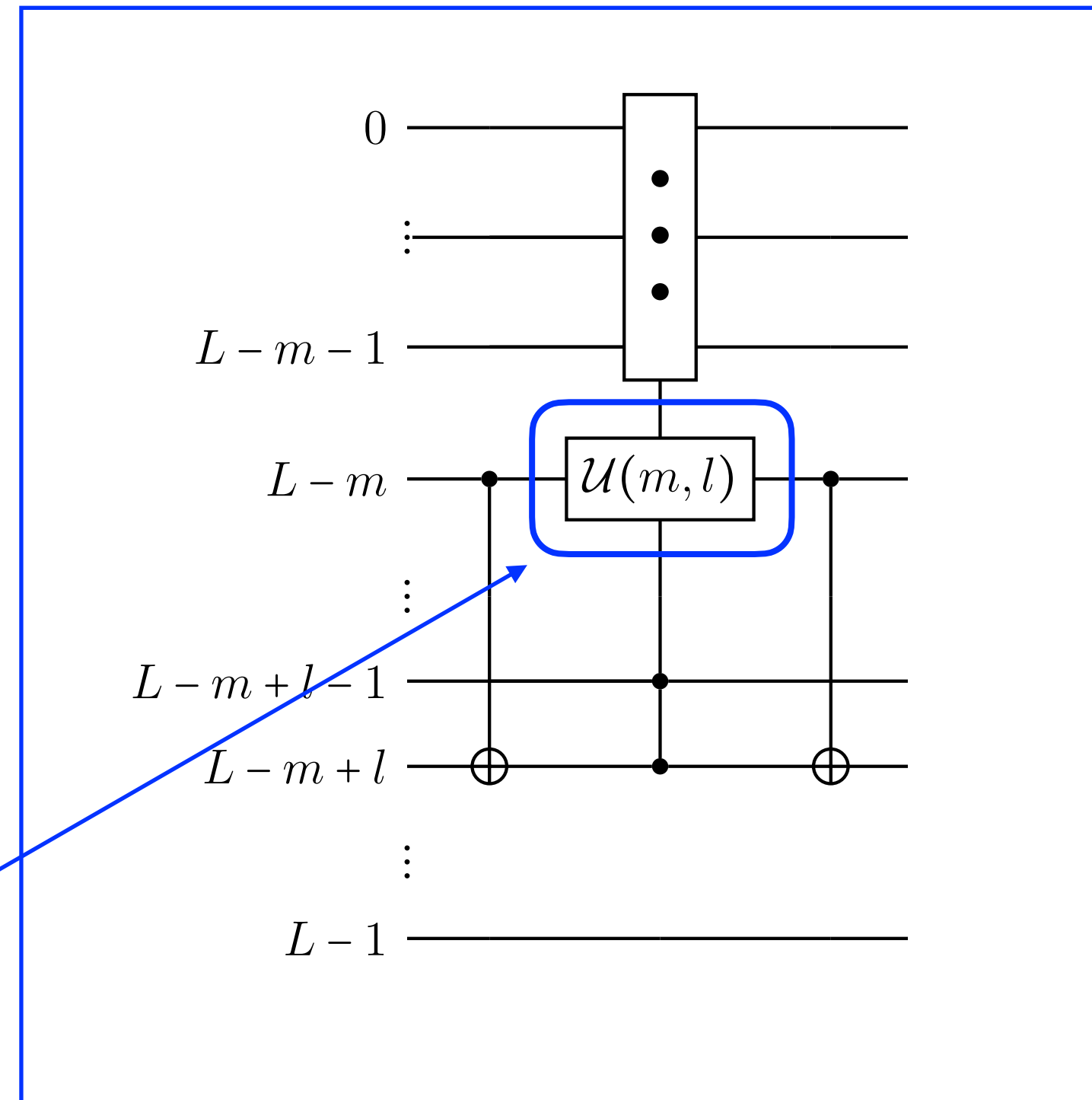
- perform  $W_m |0\rangle^{\otimes(m-l)} |1\rangle^{\otimes l} |b\rangle_{L-m}$  for fixed  $l$  for all  $|b\rangle_{L-m}$
- ordered so as to not interfere with each other

Then

$$W_m = \prod_{l=1}^{\overset{\curvearrowright}{m-1}} I_{m,l}$$

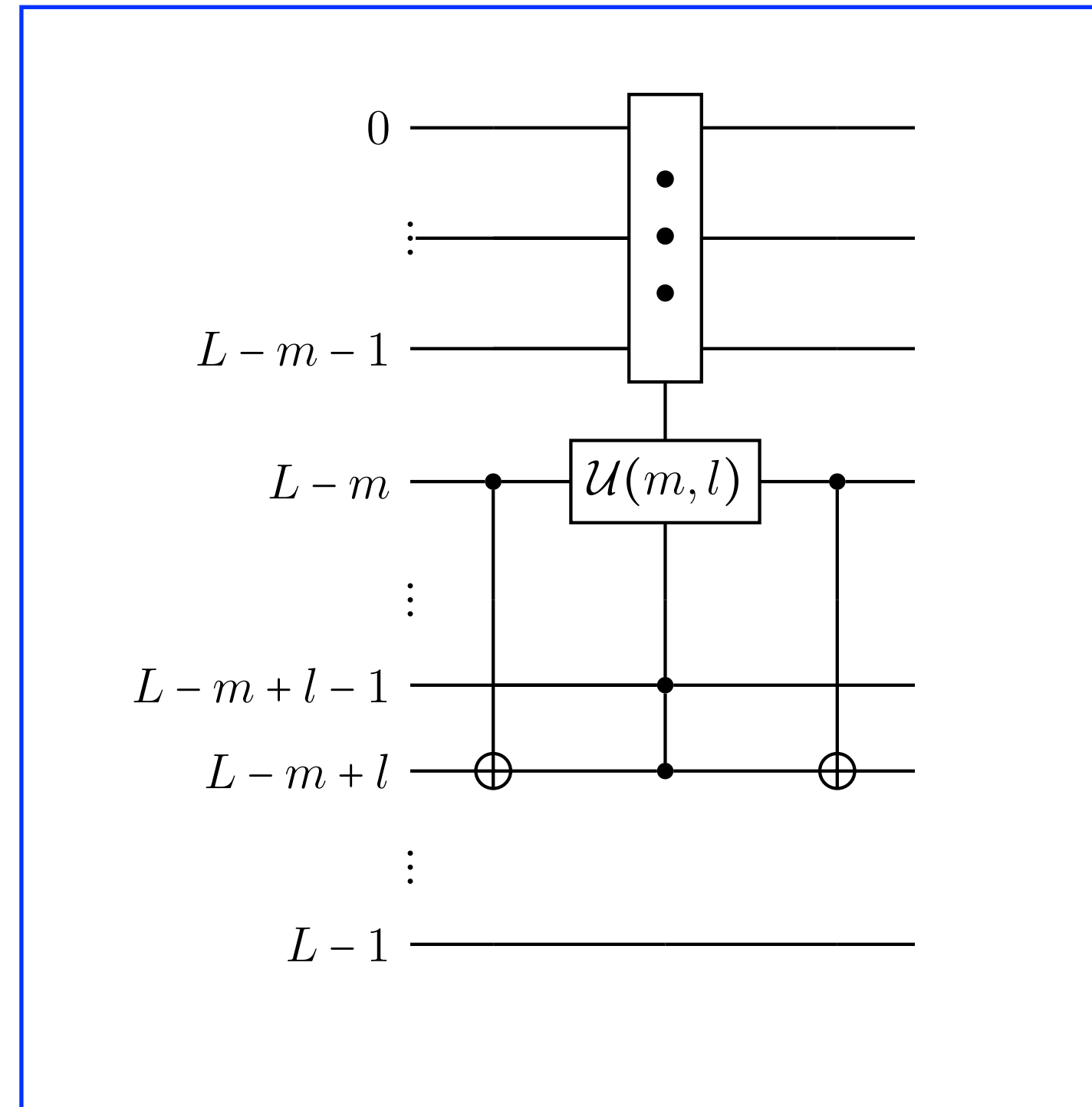
Suffices to construct  $I_{m,l}$ 's !

$I_{m,l} :$



$$\mathcal{U}(m,l) = \prod_{b \in P(L-m, M-l)} u(m,l,b)$$

$I_{m,l} :$



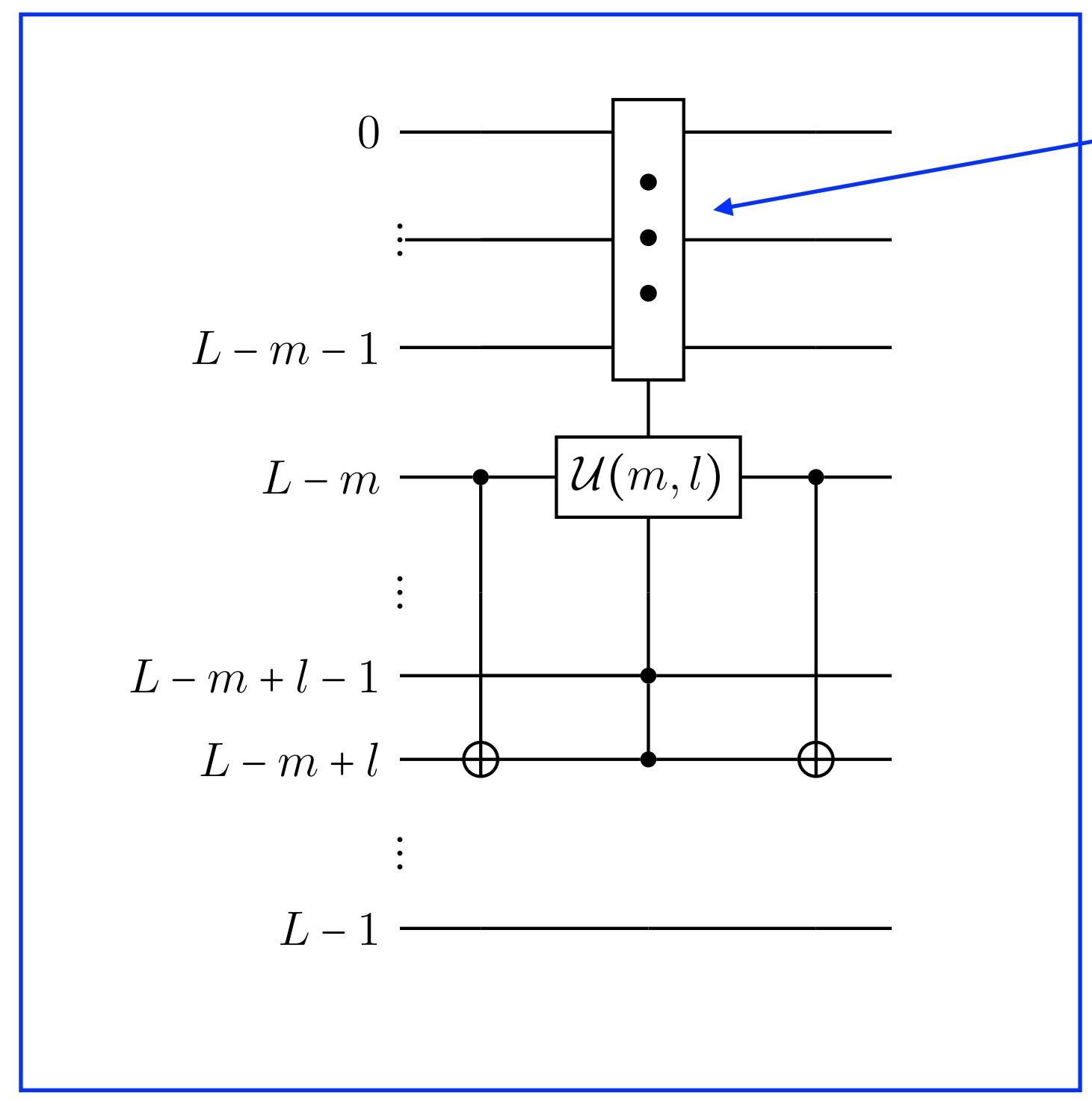
$$\mathcal{U}(m,l) = \prod_{b \in P(L-m, M-l)} u(m,l,b)$$



all possible b's for given m, l



$I_{m,l} :$



controls where b's have 1's

$$U(m,l) = \prod_{b \in P(L-m, M-l)} u(m,l,b)$$

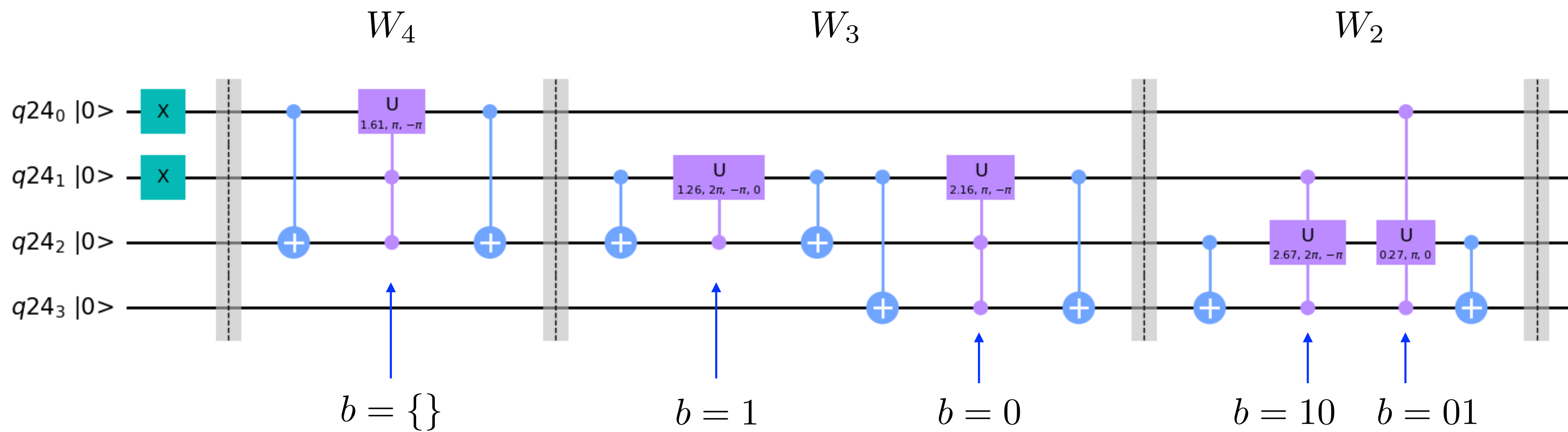
$$u(m,l,b) = \begin{pmatrix} \cos(\frac{\theta}{2}) & -e^{i\lambda} \sin(\frac{\theta}{2}) \\ e^{i\phi} \sin(\frac{\theta}{2}) & e^{i(\phi+\lambda)} \cos(\frac{\theta}{2}) \end{pmatrix}$$

$$\theta = 2 \arccos(|G(1b)|), \quad \lambda = \arg(G(0b)) - \pi, \quad \phi = \arg(G(1b)) - \lambda$$

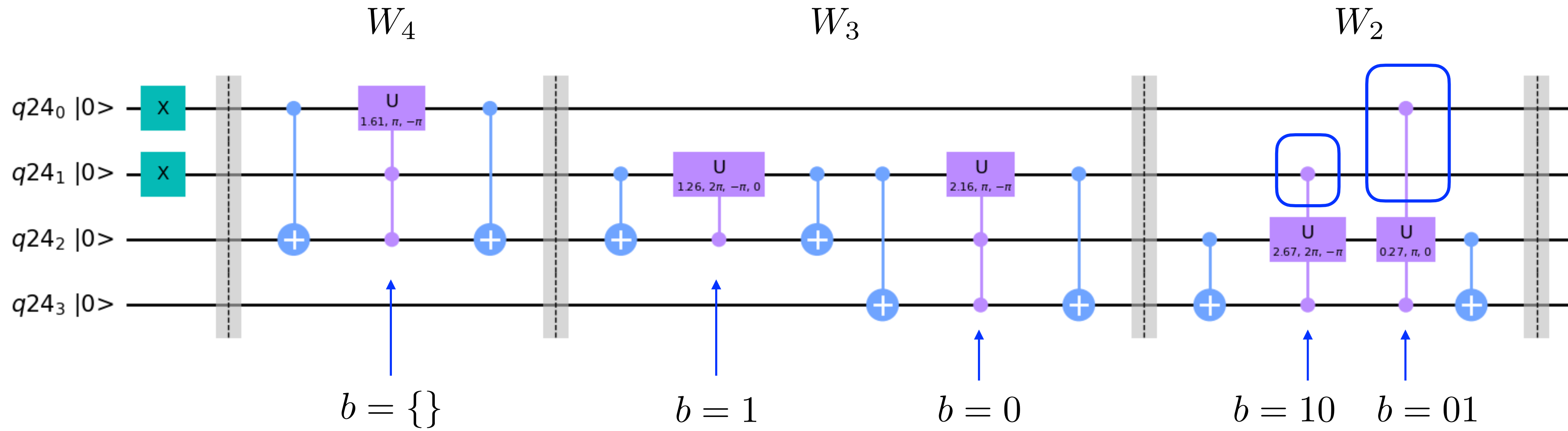
all possible b's for given m, l



Ex: (L,M) = (4,2)

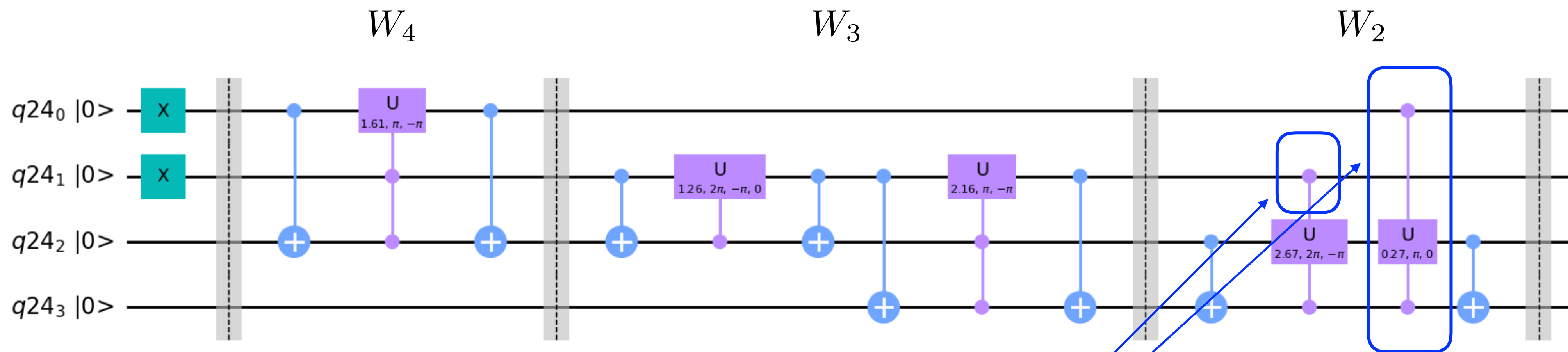


Ex:  $(L,M) = (4,2)$



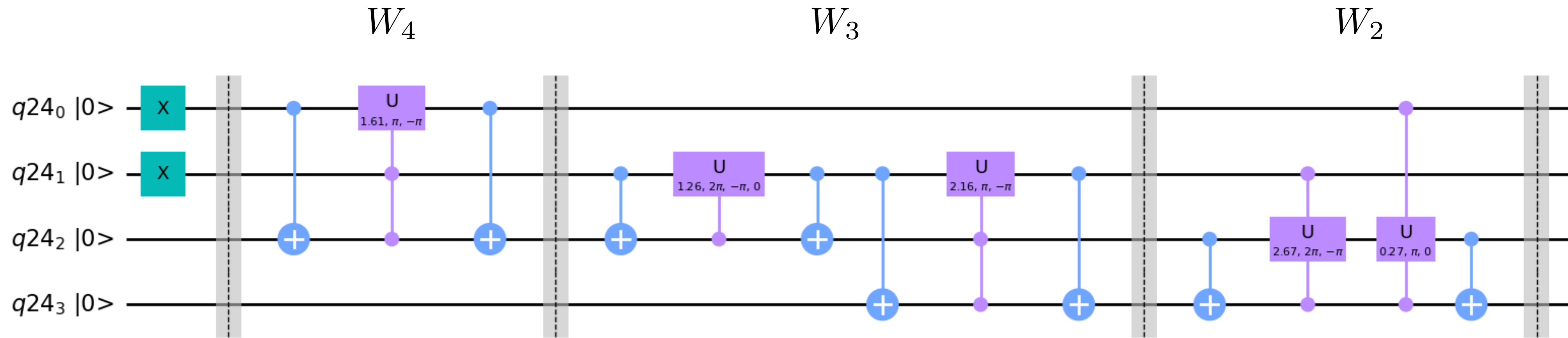
controls differentiate different b's

Ex: (L,M) = (4,2)



Same as circuit for Dicke state  $|D_2^4\rangle$  except for angles & these gates

Ex: (L,M) = (4,2)



circuit size =  $\mathcal{O}\left(\binom{L}{M}\right) \sim \# f(w)$ 's  $\therefore$  expect  $\sim$  optimal

For  $M=L/2$   $\binom{L}{L/2} \sim \frac{2^L}{\sqrt{\pi L/2}}$  😞

## 4. Bethe roots from VQE

Bethe equations are generally hard to solve. Can quantum computers help?

Variational Quantum Eigensolver (VQE): hybrid quantum/classical algorithm for estimating the ground-state energy  $E_0$  of a Hamiltonian  $\mathcal{H}$  using the variational theorem

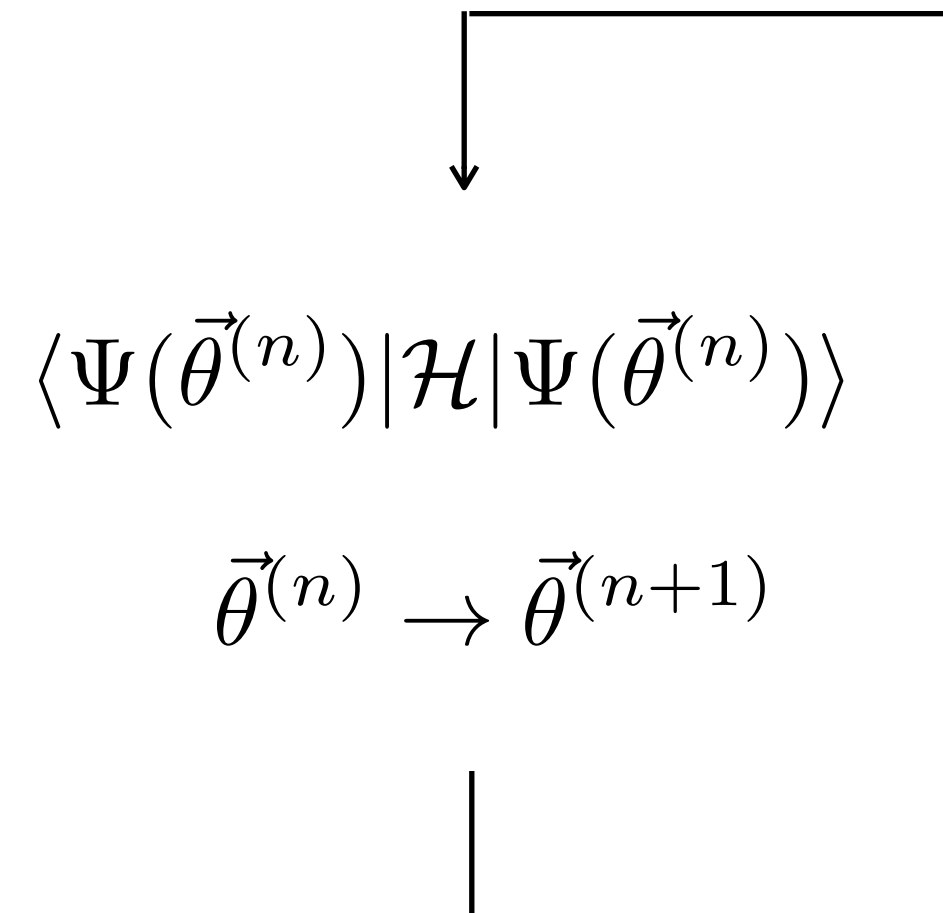
normalized trial state  $|\Psi(\vec{\theta})\rangle$        $\vec{\theta}$  parameters

iteration:  $\vec{\theta}^{(0)}$

classical simulators

~~quantum~~

classical



variational theorem  $\Rightarrow \langle \mathcal{H} \rangle \geq E_0$

Estimate Bethe roots using VQE, taking exact Bethe states as trial states, and treating Bethe roots  $\vec{k}$  as variational parameters

Raveh, RN 2404.18244

To test this idea, we instead used classical simulators.

## Classical simulators:

- **Qiskit Statevector simulator:** performs matrix arithmetic to compute exact expectation values
- **Qiskit Aer simulator:** noiseless simulation using 10,000 shots (trials)

## XXZ models:

- Closed chain

$$\mathcal{H} = \frac{1}{4} \sum_{n=1}^L \left( \sigma_n^x \sigma_{n+1}^x + \sigma_n^y \sigma_{n+1}^y + \Delta \sigma_n^z \sigma_{n+1}^z \right), \quad \vec{\sigma}_{L+1} = \vec{\sigma}_1$$

- Open chain

$$\mathcal{H} = \frac{1}{4} \sum_{n=1}^{L-1} \left( \sigma_n^x \sigma_{n+1}^x + \sigma_n^y \sigma_{n+1}^y + \Delta \sigma_n^z \sigma_{n+1}^z \right) + \frac{1}{4} \left( h \sigma_1^z + h' \sigma_L^z \right)$$

boundary magnetic fields

The state-preparation algorithm works for *any* U(1)-invariant model, including open XXZ

[Alcaraz, Barber, Batchelor, Baxter, Quispel (1987)]



## Ground-state Bethe roots:

- Closed chain  $\Delta = 2$  [Newton's method - Mathematica](#)

$L$	$M$	Energy	True roots
2	1	-2	3.14159
4	2	-2.73205	$\pm 1.94553$
6	3	-3.85577	$\pm 1.49862, 3.14159$

[Bethe roots real](#)

## Ground-state Bethe roots:

- Closed chain  $\Delta = 2$

$L$	$M$	Energy	True roots	Statevector roots	Aer roots
2	1	-2	3.14159	3.1415	3.1487
4	2	-2.73205	$\pm 1.94553$	1.9455, -1.9455	1.9623, -1.9503
6	3	-3.85577	$\pm 1.49862, 3.14159$	1.4986, -1.4986, 3.1416	1.5477, -1.4830, 3.1796

Bethe roots real

## Ground-state Bethe roots:

- Closed chain  $\Delta = 2$

$L$	$M$	Energy	True roots	Statevector roots	Aer roots
2	1	-2	3.14159	3.1415	3.1487
4	2	-2.73205	$\pm 1.94553$	1.9455, -1.9455	1.9623, -1.9503
6	3	-3.85577	$\pm 1.49862, 3.14159$	1.4986, -1.4986, 3.1416	1.5477, -1.4830, 3.1796

Bethe roots real

- Open chain  $\Delta = 1/2, h = 3, h' = 3/10$

$L$	$M$	Energy	True roots	Statevector roots	Aer roots
2	1	-0.965015	$3.14159 + 0.882174i$	$3.1415 + 0.8822i$	$3.0298 + 0.8785i$
3	2	-1.49506	$3.14159 + 0.908996i,$ 1.69883	$3.1412 + 0.9071i,$ 1.6988	$3.1537 + 0.9014i,$ 1.6930
4	2	-1.76803	$3.14159 + 0.91503i,$ 2.11689	$3.1412 + 0.9142i,$ 2.1169	$3.0944 + 0.9929i,$ 2.1473
5	3	-2.22762	$3.14159 + 0.916011i,$ 1.49569, 2.31576	$3.1399 + 0.9144i,$ 1.4958, 2.3157	$3.3120 + 1.5000i,$ 1.5449, 2.3291
6	3	-2.53682	$3.14159 + 0.916239i,$ 1.82675, 2.47141	$3.1419 + 0.9131i,$ 1.8266, 2.4712	$3.0864 + 0.9053i,$ 2.0142, 2.3022

One complex Bethe root

## Excited-state Bethe roots:

Instead of minimizing  $\langle \mathcal{H} \rangle$ , we now minimize the *variance*

[Zhang, Chen, Yuan, Yin (2020)]

$$\langle (\mathcal{H} - \langle \mathcal{H} \rangle)^2 \rangle = \langle \mathcal{H}^2 \rangle - \langle \mathcal{H} \rangle^2 \geq 0 \quad (= 0 \text{ for exact eigenstate})$$

- Closed chain  $\Delta = 2$

(selected)

$L$	$M$	Energy	True roots	Statevector roots	Aer roots
2	1	0	0	0	0.0019
3	1	-1	2.0944	2.0943	2.0938
4	2	0.732051	$\pm 0.831443i$	$\pm 0.8314i$	$\pm 0.8557i$
5	2	0.716341	$0.628319 \pm 0.835459i$	$0.6276 \pm 0.8349i$	$0.6263 \pm 0.8858i$
6	2	-1.75395	1.37766, 2.81114	1.3776, 2.8109	1.3998, 2.8293
6	3	1.18614	$0.244998 \pm 1.41247i$ , 1.6044	$0.2451 \pm 1.4120i$ , 1.6023	$0.2451 \pm 1.3341i$ , 1.3152

- Open chain  $\Delta = 1/2$ ,  $h = 3$ ,  $h' = 3/10$

(selected)

$L$	$M$	Energy	True roots	Statevector roots	Aer roots
2	1	0.715015	1.30258	1.3026	1.3025
3	1	-0.869852	$3.14159 + 0.911371i$	$3.1417 + 0.9113i$	$3.1683 + 0.8783i$
4	2	-0.224189	$3.14159 + 0.916221i$ , 0.2264i	$3.1413 + 0.9164i$ , 0.2284i	$3.0311 + 0.9359i$ , 0.2148i
4	3	-0.128194	$3.14159 + 0.916237i$ , 0.93789, 0.245389i	$3.1416 + 0.9174i$ , 0.9382, 0.2474i	$3.0528 + 0.9138i$ , 0.9046, 0.2146i
5	4	-1.61607	$3.14159 + 0.916185i$ , 0.514675, 1.16211, 2.43263	$3.1421 + 0.9027i$ , 0.5044, 1.1618, 2.4332	$3.0030 + 0.7578i$ , 0.4451, 1.0290, 2.4882
6	5	0.21968	$3.14159 + 0.916291i$ , 0.667057, 0.32195i, 1.12044 $\pm$ 0.160175i	$3.1412 + 0.9156i$ , 0.6252, 0.3029i, 1.1347 $\pm$ 0.1956i	$3.1218 + 0.7523i$ , 0.5055, 0.5615i, 1.0101 $\pm$ 0.2531i

multiple complex Bethe roots

## 5. Outlook

- Deterministic algorithm for preparing exact Bethe states

Works for *any*  $U(1)$ -eigenstate

- Used for estimating Bethe roots

- Exploit integrability?

- Reduce circuit size and/or depth? ancillas, measurements, feedforward operations Piroli, Styliaris, Cirac 2024

Mao, Tian, Sun 2024

- Higher spin? Higher rank?

- Models without  $U(1)$  symmetry?

Thank you for your attention!